



# **MICROPROCESSOR SYSTEMS (IAS0430)**

Department of Computer Systems  
Tallinn University of Technology

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## KERNEL VS USER

- Now that we know what makes a CPU.
  - What is the CPU made of?
  - How does it run?
  - What makes it operate?
  - How does it know what to operate on?
  - Where does it find things?
- Next is to learn how it interacts with **programs**:
  - More specifically, **which program is allowed to do what?**
  - In reality, the CPU must protect itself and its operations.
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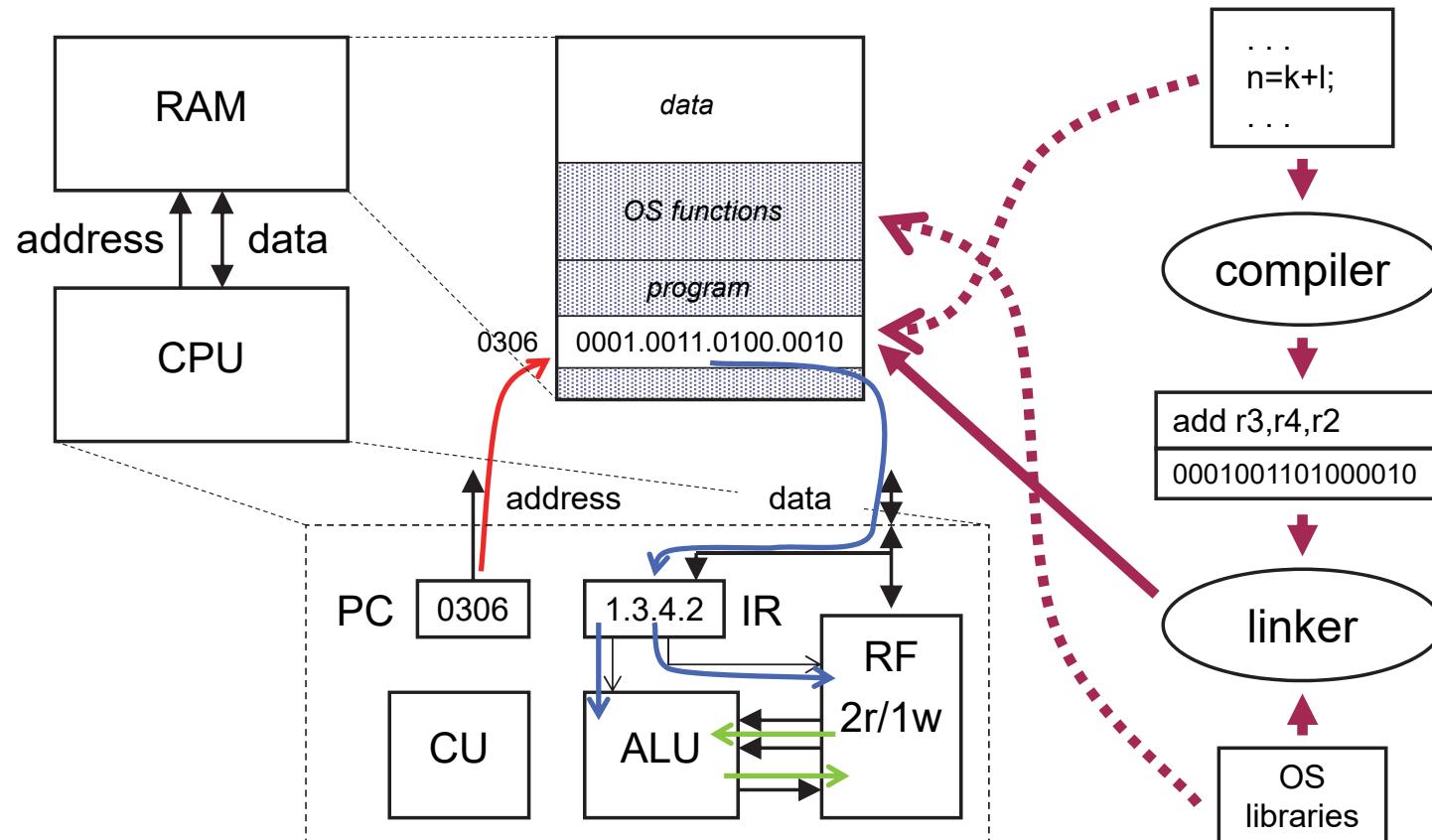
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      - **Why?**
      - Some location in memory **contain sensitive information** that must not be changed/updated/ altered/removed – **NOT A TOY TO PLAY WITH**

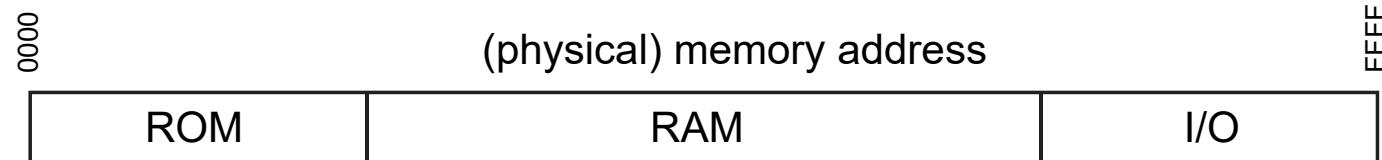
# RUNNING A PROGRAM



# PROGRAM IN MEMORY

- **Single** application/program
  - Memory layout
    - Program – application + system
    - Data – static / dynamic / stack
  - Program running
    - Starting up (booting)
      - Initializing processor
      - Initializing program
    - Program in loop
    - System calls

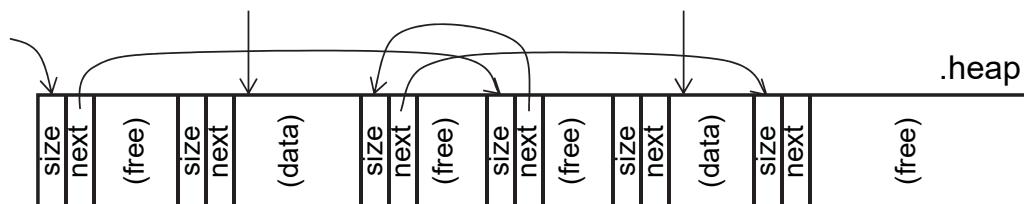
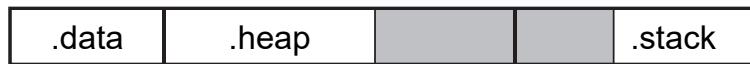
## MEMORY LAYOUT – SINGLE APPLICATION



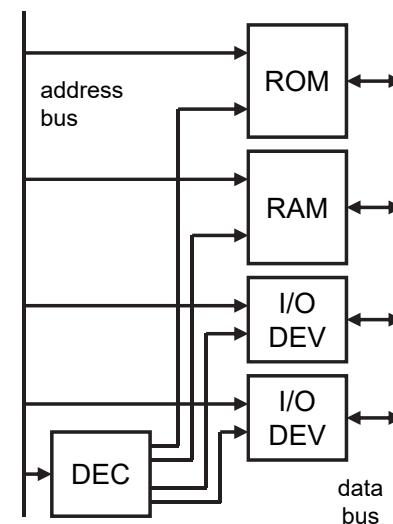
ROM



RAM



The physical module  
is selected/activated  
by decoder(s)

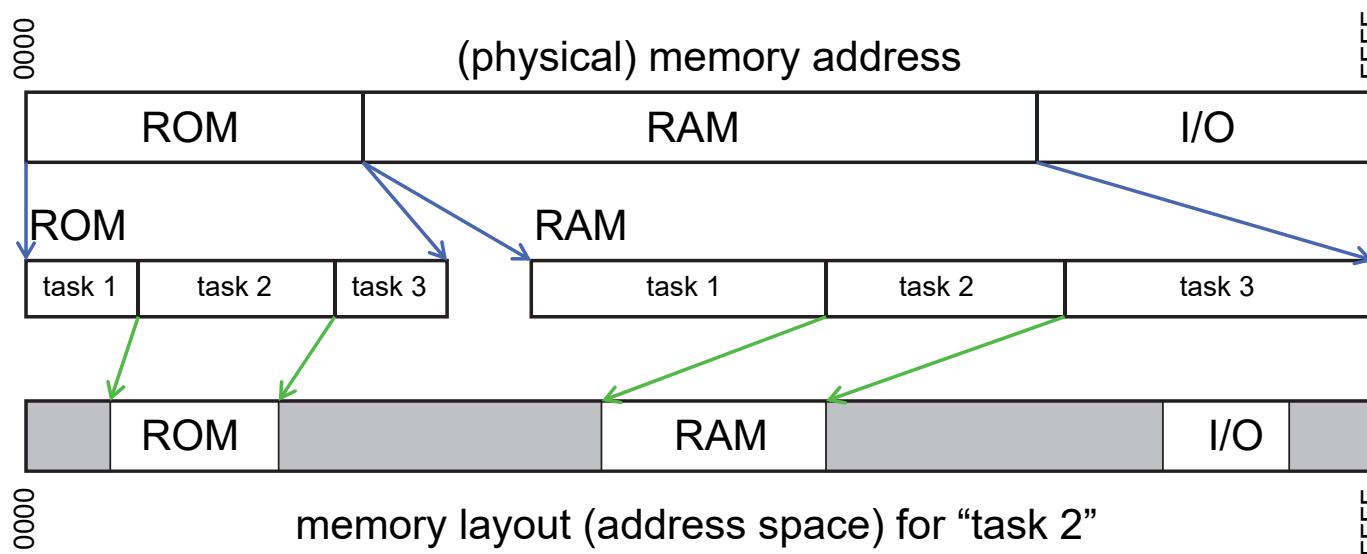


# PROGRAMS IN MEMORY

- **Multiple** applications/programs
  - in addition to the single application
- Memory layout
  - Different programs & their data at different memory locations
  - Protecting data
    - MMU – Memory Management Unit
- Programs running
  - Initializing system
  - Changing applications/programs/tasks
    - scheduling, saving CPU status, ...

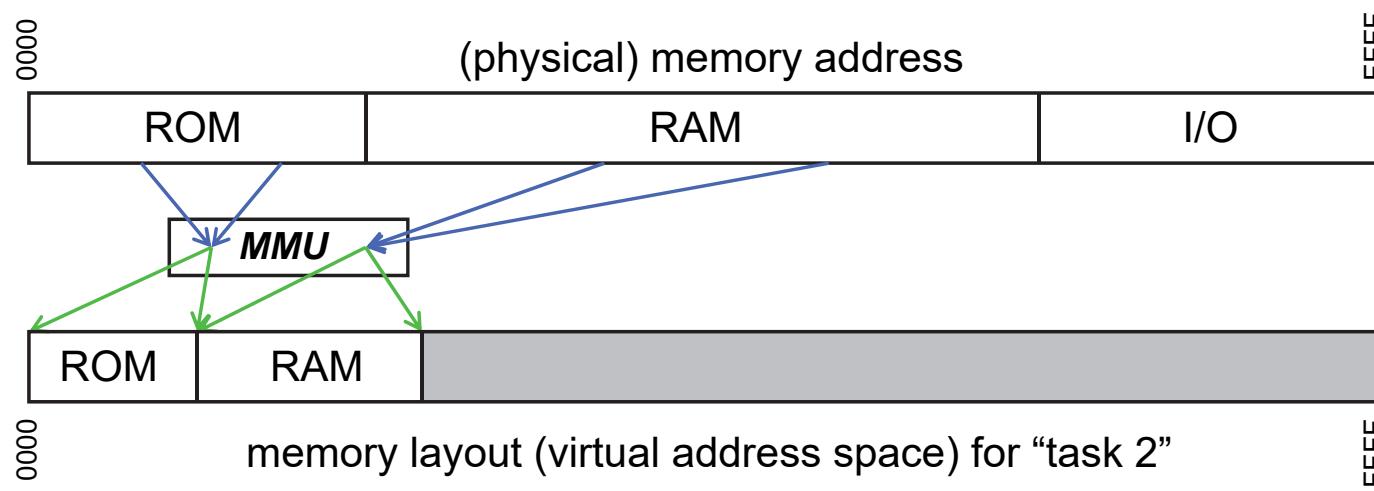
## MEMORY LAYOUT – MULTIPLE APPLICATIONS

- **Every** application/program at different physical address
  - Internal memory layout is the same for every application (.data, .heap, etc.)
  - MMU is used to protect from wrong/faulty memory accesses
  - Linking is complicated – addresses must be defined separately



## MEMORY LAYOUT – MULTIPLE APPLICATIONS

- Memory layout is the **same** for **all** applications/programs
  - virtual address space (virtual memory)
- Internal memory layout is the same for every application (.data, .heap, etc.)
- MMU is used to protect from wrong/faulty memory accesses and also to “translate” the virtual address into physical address
- Linking is simple – same for all applications/programs



## KERNEL VS USER

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    - We will not talk about the **OS** today, that's for next week.
  - Today we will talk about what makes the **OS** so ..... **OSy**
- Today we will talk about **Kernel** and **User** Modes
- A **Kernel** is the soft part of a nut – **Yummy**
  - But what we care about is that it is **contained within a shell**
- That is why we call the core of the computer – a Kernel
  - What is the shell of a **Computer Kernel?**

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  - What is the shell of a **Computer Kernel?**
  - The **OS**

## KERNEL VS USER

- **Kernel mode:**

- Also known as the **super user** mode
- The CPU is full on operation work. All is granted and all can be done.
- When a program runs in Kernel mode
  - It means that the program is being executed with the CPU is unrestricted!
  - It has complete access to the **hardware** (CPU, ports, memory, etc.)
    - Can perform any operation allowed by the architecture.
    - Can signal any I/O device.
  - It can execute any **CPU instruction** and reference any **memory location** that it may need.
- Programs that run in Kernel mode are usually the most **trusted and essential functions and services** of the OS.
- Kernel mode crashes (programs that fail to execute properly) are catastrophic.
  - Will almost always result in complete system halt
  - Example of this happens in Windows?

# KERNEL

## ■ Kernel

- Also known as the operating system's core.
- The part of the OS that runs in memory.
- Where the OS interacts with hardware.

A problem has been detected and Windows has been shut down to prevent damage to your computer.

THREAD\_STUCK\_IN\_DEVICE\_DRIVER

If this is the first time you've seen this Stop error screen, restart your computer. If this screen appears again, follow these steps:

Check to make sure any new hardware or software is properly installed. If this is a new installation, ask your hardware or software manufacturer for any Windows updates you might need.

If problems continue, disable or remove any newly installed hardware or software. Disable BIOS memory options such as caching or shadowing. If you need to use Safe Mode to remove or disable components, restart your computer, press F8 to select Advanced Startup options, and then select Safe Mode.

- Programs and functions.

Technical information:

- Kernel errors.

\*\*\* STOP: 0x0000000EA (0x00000000, 0x00000000)

- 
- 
- 

- The blue screen of death

## KERNEL VS USER

### ▪ User Mode:

- When a program runs in user mode, also known as **slave mode**...
  - It has no direct communication with the **hardware** (CPU, ports, memory, etc.)
  - It can only reference **memory locations** relevant to its own execution (its own memory space).
  - It must communicate with the system **Application Program Interface (API)** in order to access hardware and memory.
- Programs that run in User mode are usually software the interact with the user of the system directly (power point, chrome, etc..)
- User mode crashes are **manageable** and **easy to recover** since a user mode program has no direct access to the system services or resources. Thus, does not effect the overall all functionalities of the system.
- Kernel mode aims to **increase stability of the system by managing program access to volatile and sensitive resources**.

## KERNEL VS USER

- **So to summarize:**
- **Kernel Mode**
  - Only specific system services and functions can run on CPU Kernel mode
    - Because they are important for performance
    - And they make little threat to system resources
    - Are not controlled by the User!
  - Kernel mode helps protect the core of the system functionality.
- **User Mode**
  - Programs can only access resources allocated to them by the system.
    - This is for keeping the system stable
    - Prevents programs from accessing hardware that does not belong to them
    - Allow interaction between the system hardware and the User without putting the system in jeopardy.

## KERNEL VS USER

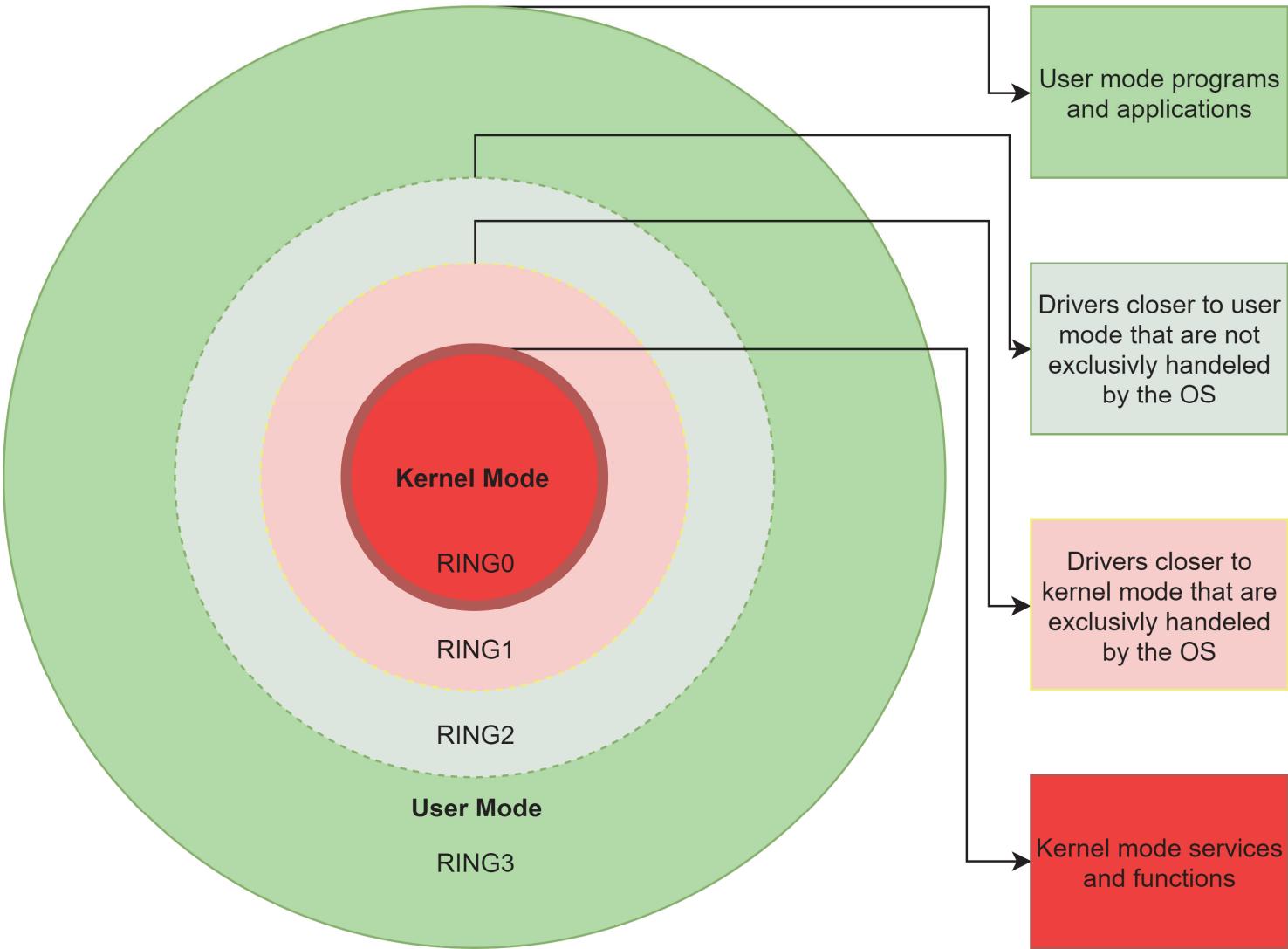
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## KERNEL VS USER

- **But**, what about when user mode needs to access hardware that is not managed by the system?
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- **RING0**

Exclusive Kernel Mode programs. **Hardware management software and device controllers communication software.**

This also includes **OS services** that need direct access to hardware

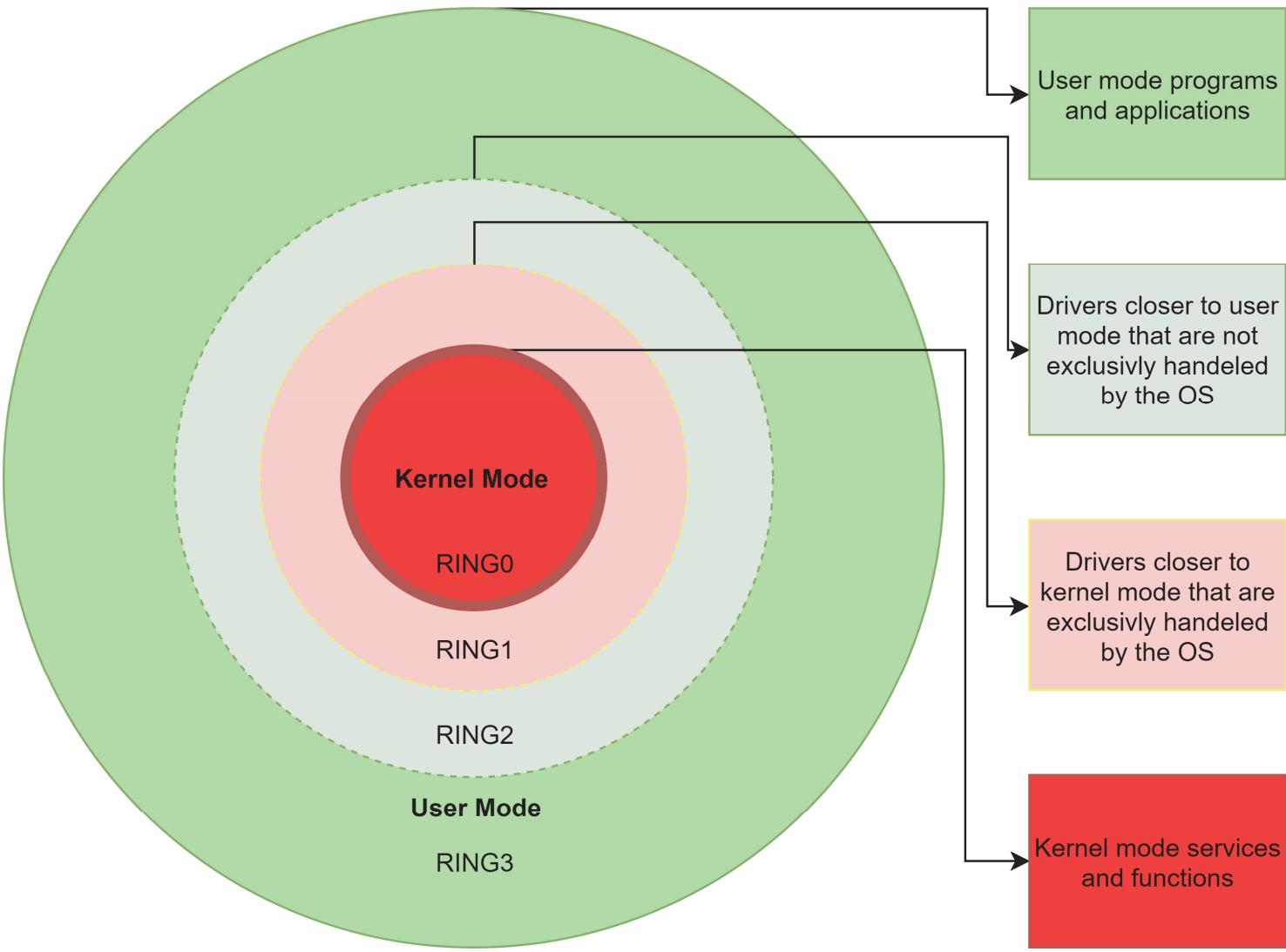


## KERNEL VS USER

### ▪ RING1

Kernel Mode related Drivers. These are drivers such as the video card and LAN card that need driver controllers (**software**) to manage the physical drivers (**hardware**) using the OS services.

**Those drivers do not get full access to kernel mode**, but they have certain level of privilege that allows them to communicate on kernel mode-like way. This mostly managed by the OS, but also the system kernel manage them as well.

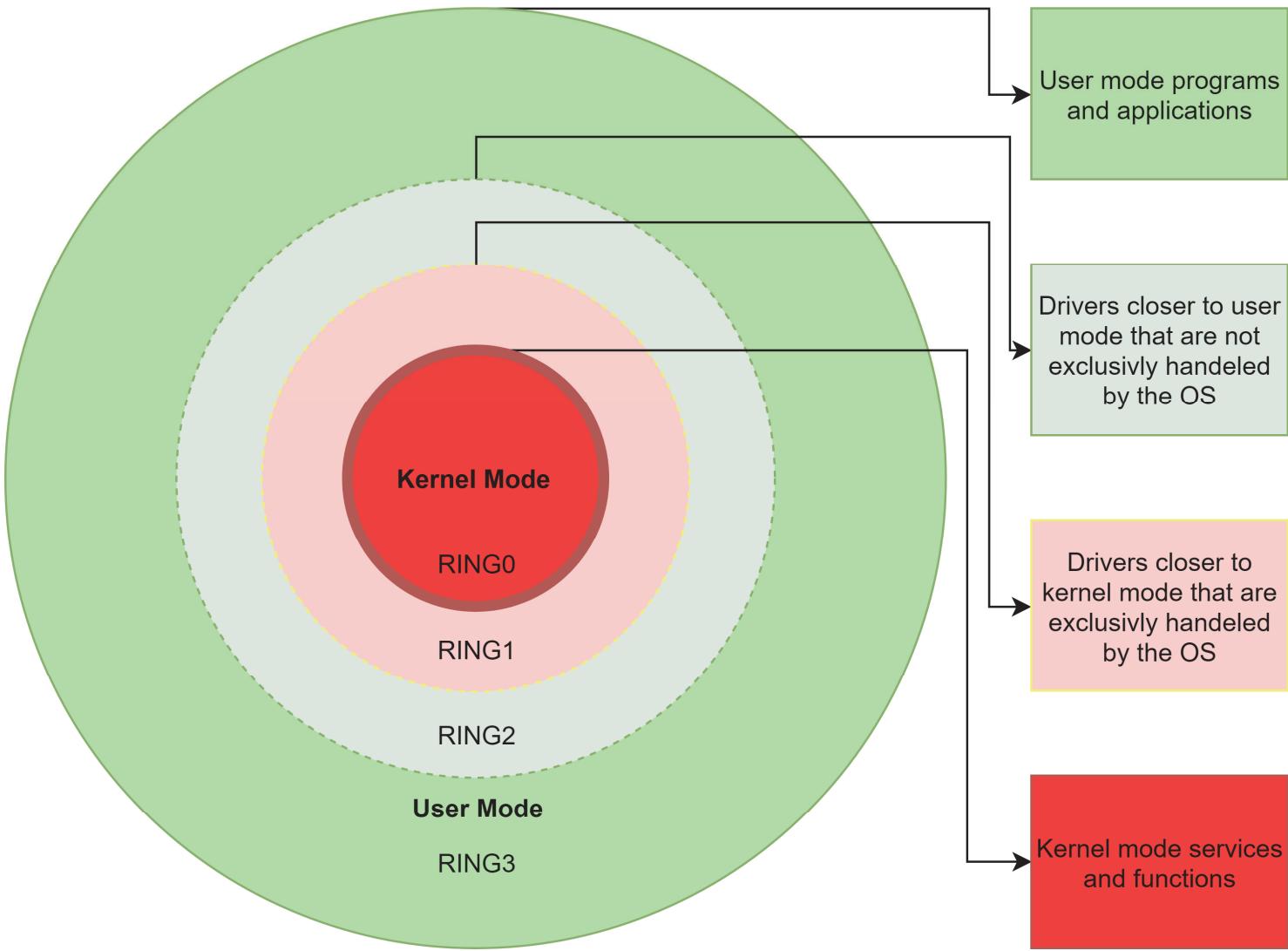


## KERNEL VS USER

### ▪ RING2

User mode related Drivers.  
Drivers that do not necessarily need OS services and can be managed by software independently than the OS. Such as **mouse**, **keyboard**, sometimes **USB**, **printers**, etc....

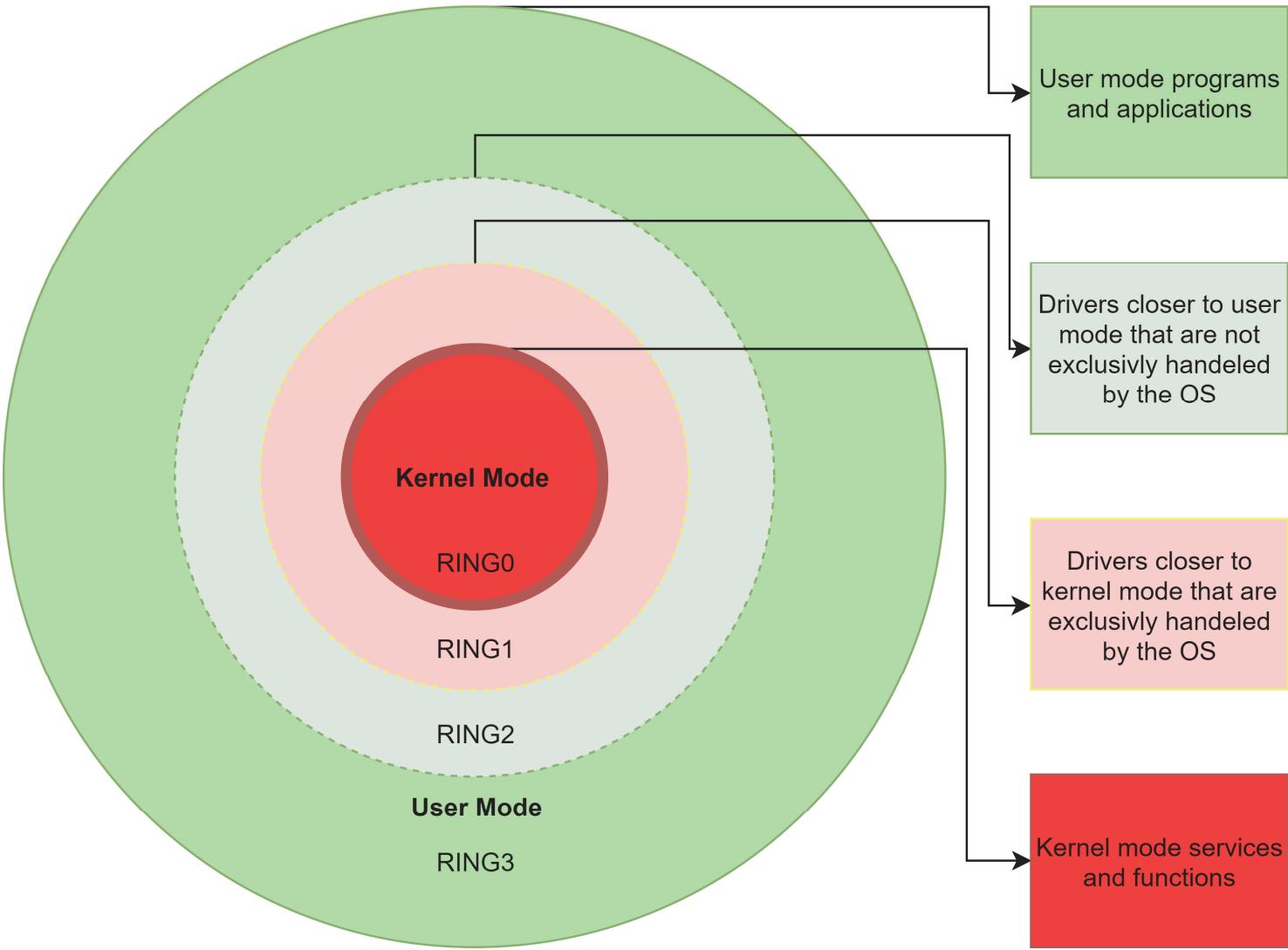
Those drivers do not get any privileges, but they will get some restricted access to kernel mode when needed. This is mostly managed by the OS.



## KERNEL VS USER

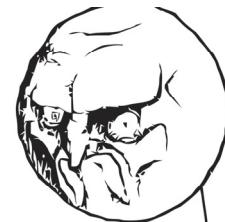
### ▪ RING3

Includes all other types of programs and applications that run on User mode. These programs must directly communicate with API if access to Hardware was needed.



## KERNEL VS USER

- **But again**, why do we need all this mess, I mean, is it not easier if we can just run things on computers without having to worry about stuff....



NO.

## WHY WE NEED THOSE MODES

- **Kernel and User modes provide features to the system**

- **Security:**

- Create layers of protection around data, user data, memory, and internal system operations
    - Preventing bad programs from interfering with system performance and OS
    - Provide a recovery protocol in case of process or program scheduling failure.
    - Different parts of the system run separately, allowing modification to be done when necessary without intervening with other parts.

- **Privileges:**

- We need clarify which programs can do what.
    - What resources can a program use.
    - Control which processes can have more authority on resources.
      - Memory
      - Registers
      - I/O devices
      - CPU time and internal system functions

## WHY WE NEED THOSE MODES

- Kernel and User modes provide features to the system
  - Control of execution flow:
    - Control how execution is handled and how programs are scheduled.
    - Kernel mode programs **can** be of higher importance
      - Thus, they must run on higher priority.
  - Giving more priority to **system calls** or **syscalls**
    - When a user mode program requires access to HW, it issues a syscall.
    - A user mode program with a syscall must be handled before other user mode programs without a syscall

## **SWITCHING**

- **Switching in between**

- Ok, we know what is kernel and user modes and we know why we need them, but how does a program go from user mode to kernel mode and the other way around.

# SWITCHING

## ▪ Switching in between

- Ok, we know what is kernel and user modes and we know why we need them, but how does a program go from user mode to kernel mode and the other way around.
- Case: access to keyboard by **notepad**
  - Normally, **the CPU runs a program like notepad in User mode**, but it also switches to kernel mode when needed.
  - When a notepad needs access to the keyboard, it reaches out to the **OS API**.
    - It sends a **trap instruction** to the OS.
      - A trap is a special function instruction that requests access to an OS service.
      - Also known as a **syscall** instruction in MIPS ISA.
    - When the OS received the **trap**, it signals the CPU to switch to Kernel Mode to run the OS service that deals with keyboard accesses – **the trap handler**.
    - Once the CPU is in kernel mode, it jumps directly to a fixed location in memory.
      - This location is part of the OS that can only be accessed in kernel mode.
      - The location is **the trap handler** program of the OS.

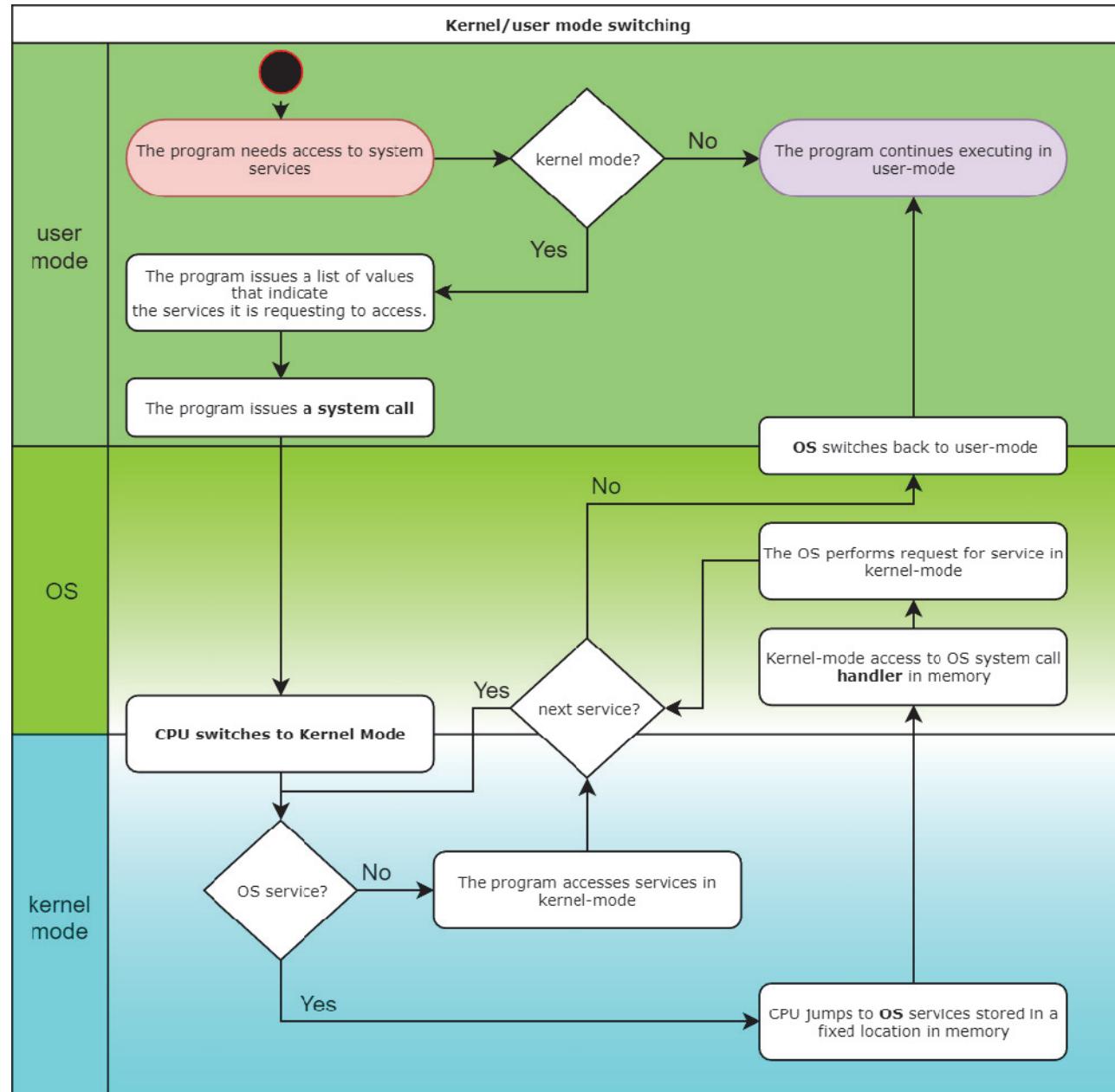
## SWITCHING

### ▪ **Switching in between**

- The trap handler starts by jumping to **the location in memory** where the keyboard device driver is – again, only accessible in kernel mode.
- Now that we are where the driver is, the **OS** (using the driver program) **accesses the keyboard on-behalf of the notepad program**.
  - It reads the input of the keyboard and sends it to the CPU (still in kernel mode) to be handled.
  - The input from the keyboard is then processed.
- Once the input from the keyboard is processed, the CPU signals the OS that the trap sent by the notepad program is fulfilled.
- The OS replies to the notepad that the trap is handled.
  - Then it switches the CPU back to User mode.
- The notepad program continues executing in user mode.
  - In reality, the notepad itself (program instructions) are never at any case are executed in kernel mode.
  - To the note pad, the trap instruction is a **magical instruction!**
    - It appears that the whole system call was done in one instruction.

# SWITCHING

- **Switching in between**



## MORE IS MORE

- **This not about kernel mode and user mode:**

- There is no single kernel mode!
  - There are a variety of modes that “kernel mode” refers to.
    - Depend on the program permissions, system functions, and OS services.
    - Same for user mode.
- Programs other than the OS rarely run in kernel mode
  - Some exceptions like anti virus programs, Device control BIOSes (Nvidia, Intel, etc.)
    - That is why those require a direct access before the OS boots.
    - Otherwise, the OS will block them from accessing kernel mode.
- Almost no program (be it OS service or BIOS service) can run on full kernel mode.
  - No program can control everything!
  - So even in kernel mode there are restrictions in some cases.